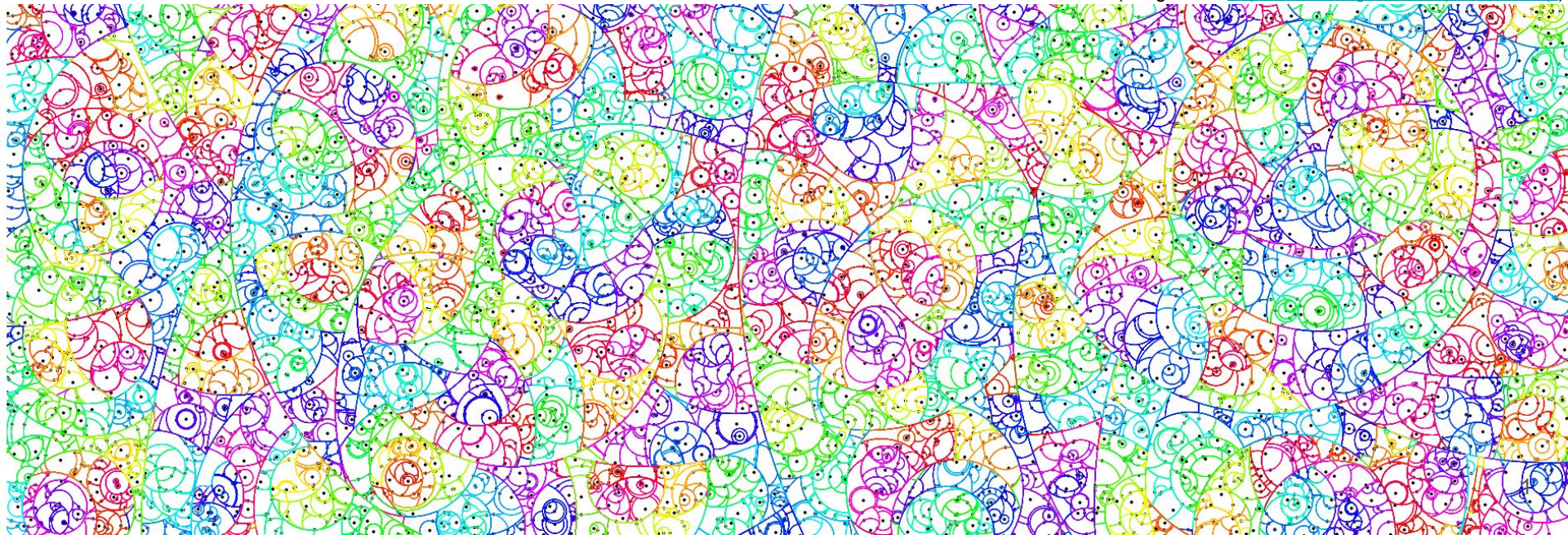


# Vantage Point Trees

Based on “Data Structures and Algorithms for Nearest Neighbor Search in General Metric Spaces”,  
Peter N. Yianilos (1993) - <http://algorithmics.lsi.upc.edu/docs/practicas/p311-yianilos.pdf>

(Image from <https://fribbels.github.io/vptree/writeup>)



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- Nearest Neighbor Search
  - given  $X$  a dataset of  $N$  points and a query point  $q$ , find the point  $x_i$  minimizing the distance between itself and  $q$ .

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- Set of possible solutions
  - Brute Force,  $O(N^2)$

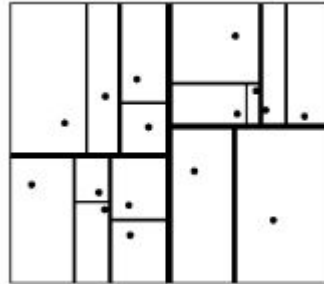
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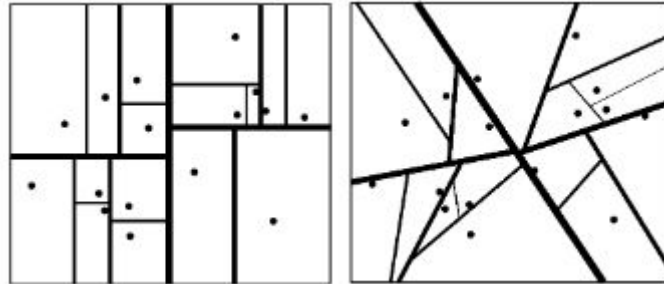
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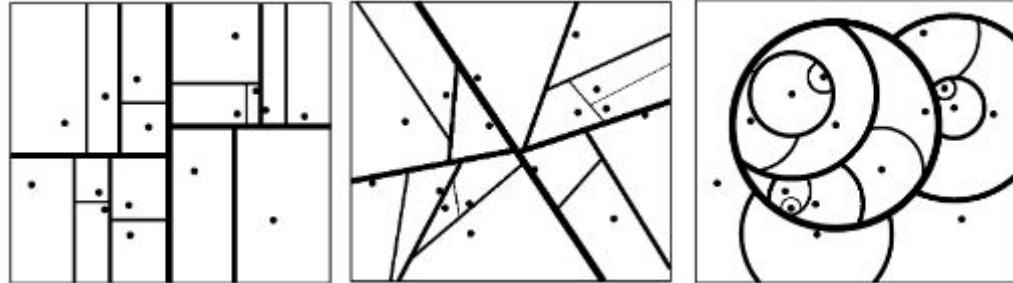
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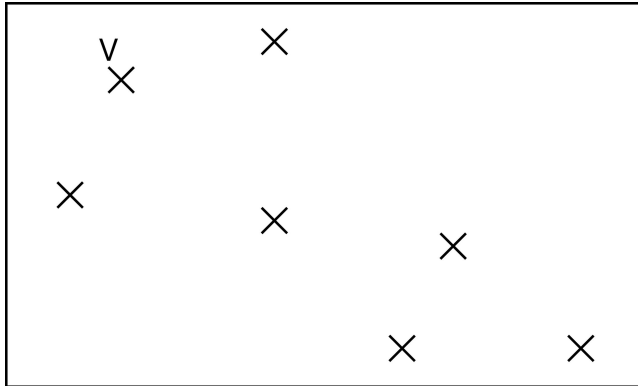
- Brute Force,  $O(N^2)$
- Metric Trees
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  - **Vantage Point Trees**
    - $O(N \log(N))$  to build it
    - $O(\log(N))$  for a NN Search



(Images from [https://link.springer.com/content/pdf/10.1007/978-3-540-88688-4\\_27.pdf](https://link.springer.com/content/pdf/10.1007/978-3-540-88688-4_27.pdf))

# Vantage Point Tree: Construction

1. Select a vantage point  $v$  in  $X$  (eg. following a uniform distribution);

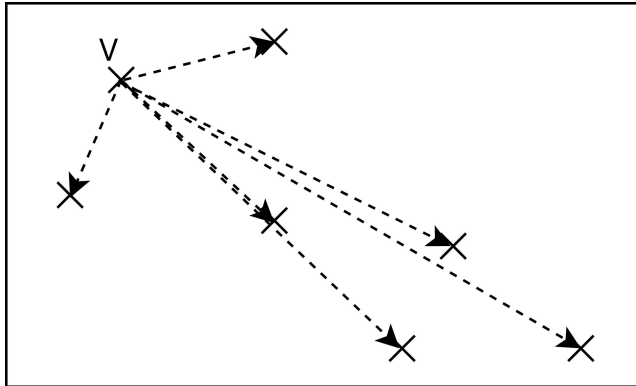


$v$

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2. Compute the distances  $d(v, x_i)$  between  $v$  and each point  $x_i$  in  $X$ ;

Euclidean in this example but works with any metric respecting the triangle inequality!

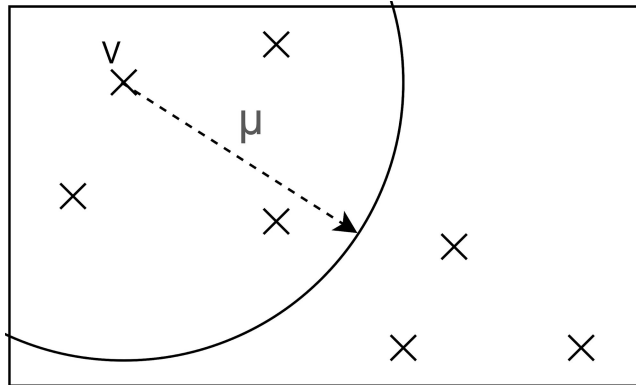


v



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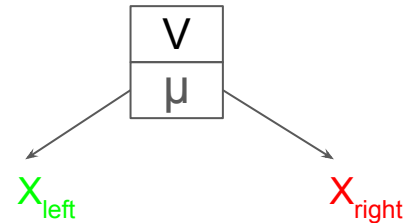
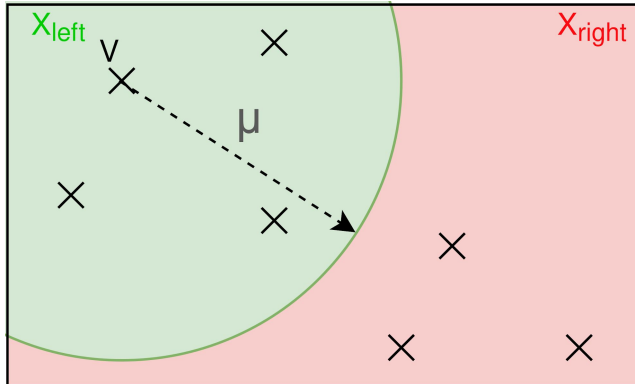
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$v$
$\mu$

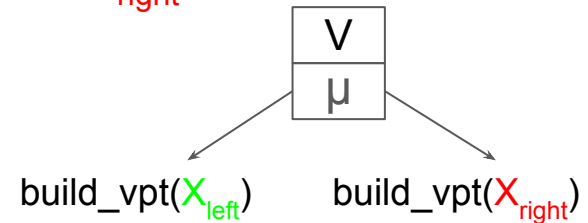
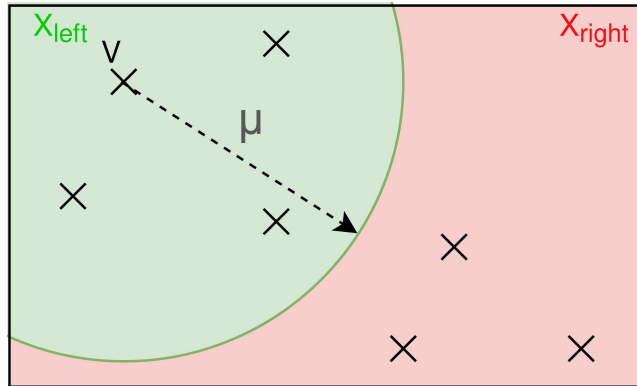
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  - a.  $X_{\text{left}}$  the set of points closest to  $v$  is put at the left
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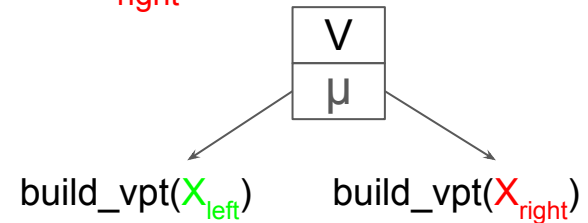
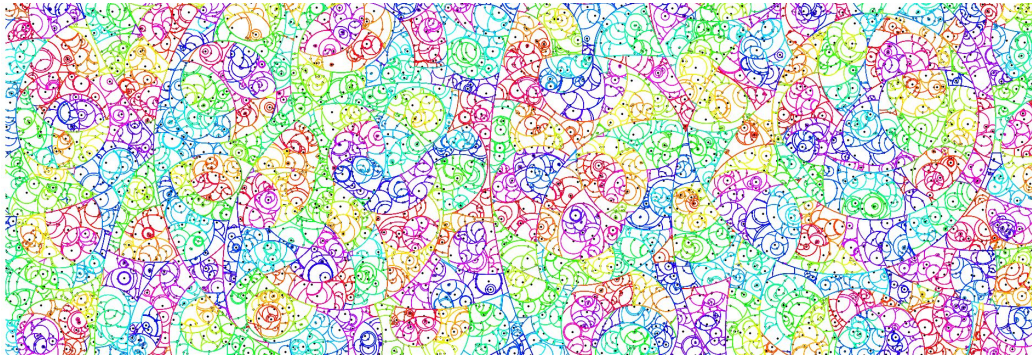
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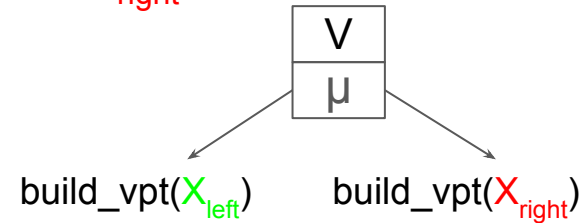
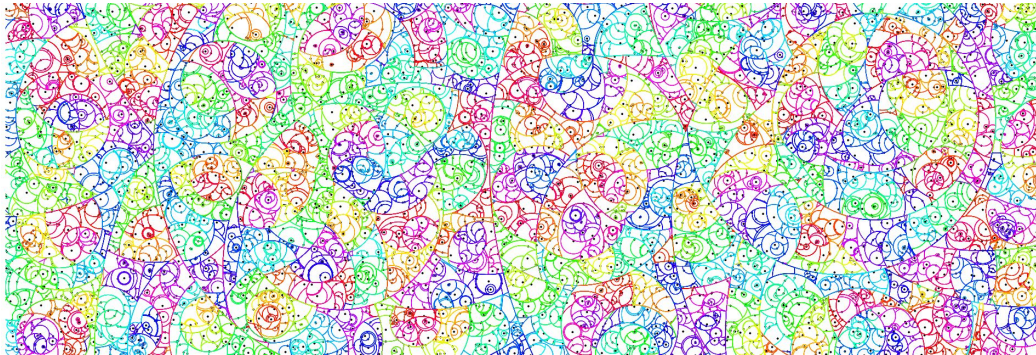
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DEMO  
<https://fribbels.github.io/vptree/writeup>




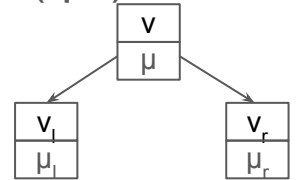
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


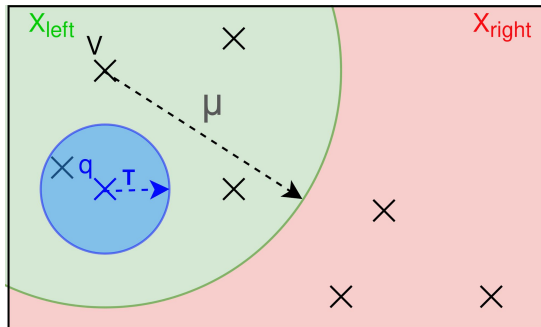
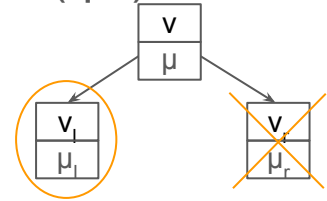
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
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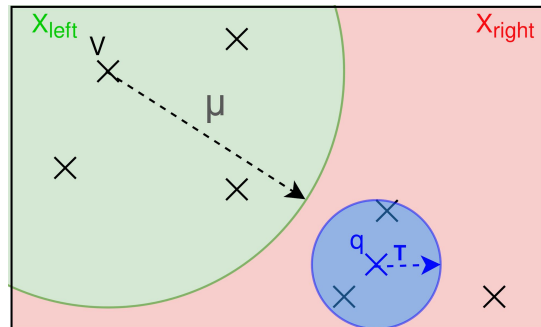
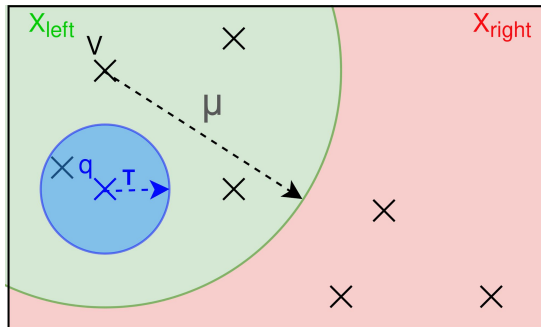
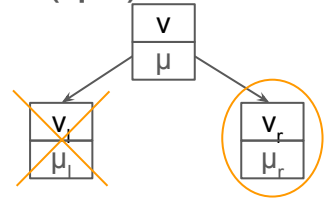
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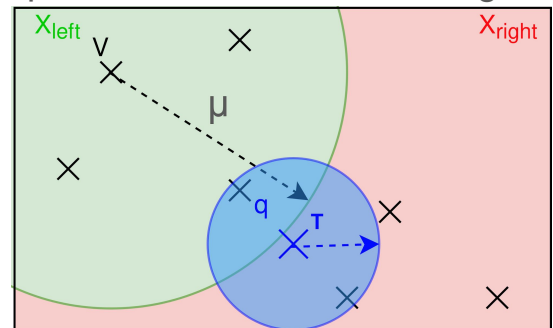
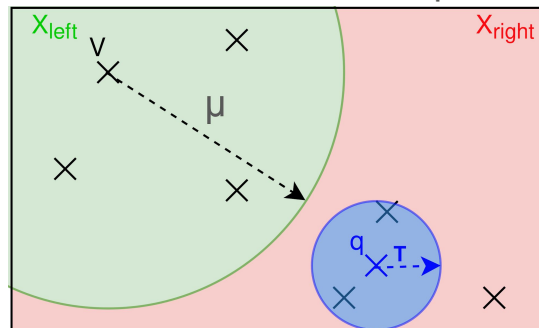
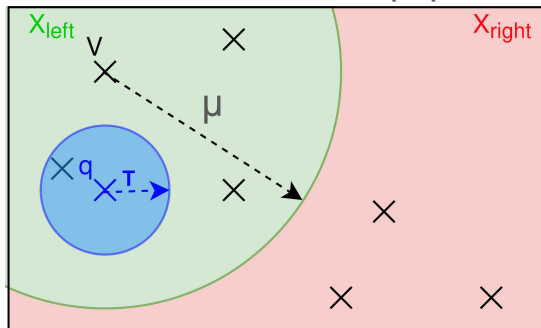
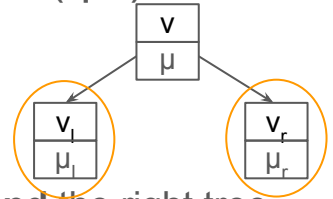
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- c. If the  $q$ -sphere is inside and outside the  $v$ -sphere: explore the left tree **and** the right tree



# Vantage Point Tree: Extensions

- $VP^S$ -Tree
  - Instead of retaining only the median, a node retains 2 values per ancestor corresponding to the boundaries of the corresponding subspace

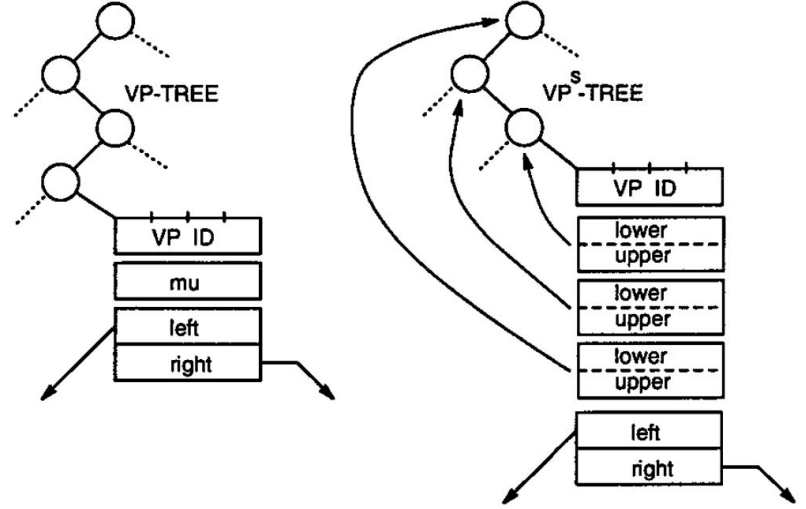


Figure 5: Sample 32-bit machine data structure implementations for the most basic vp-tree, the  $vp^S$ -tree, and the  $vp^{sb}$ -tree.

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  - Reduced memory consumption

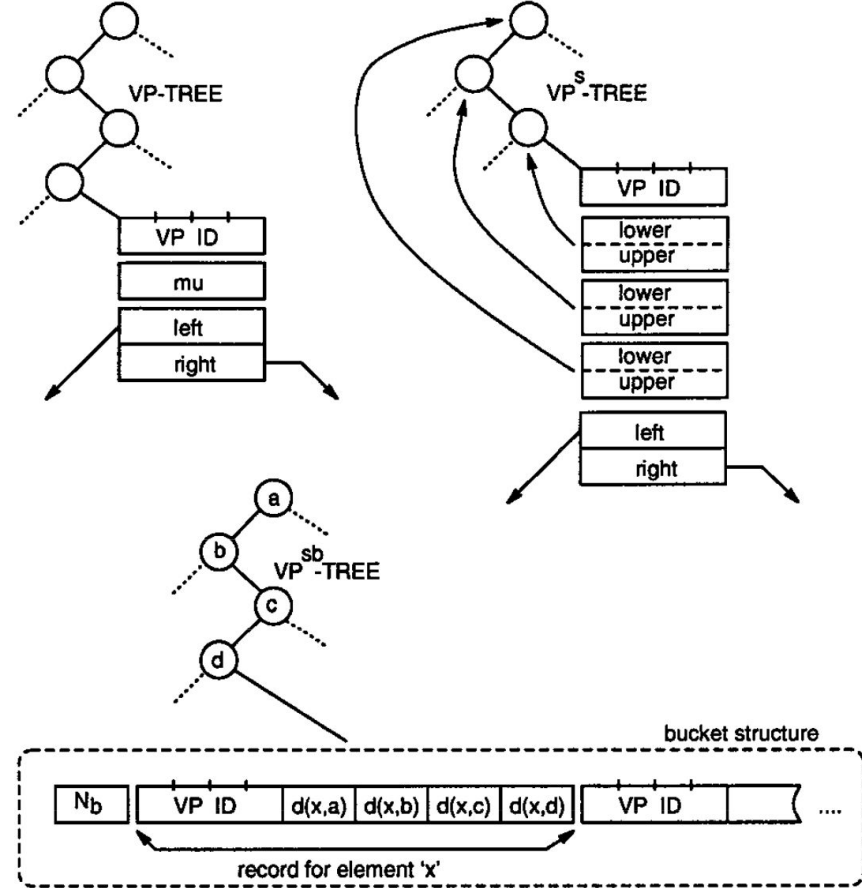


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  - Reduced memory consumption
- VP Tree in non-metric space
  - With Bregman divergences:
 

“Bregman Vantage Point Trees for Efficient Nearest Neighbor Queries”  
(Nielsen et al., 2009)

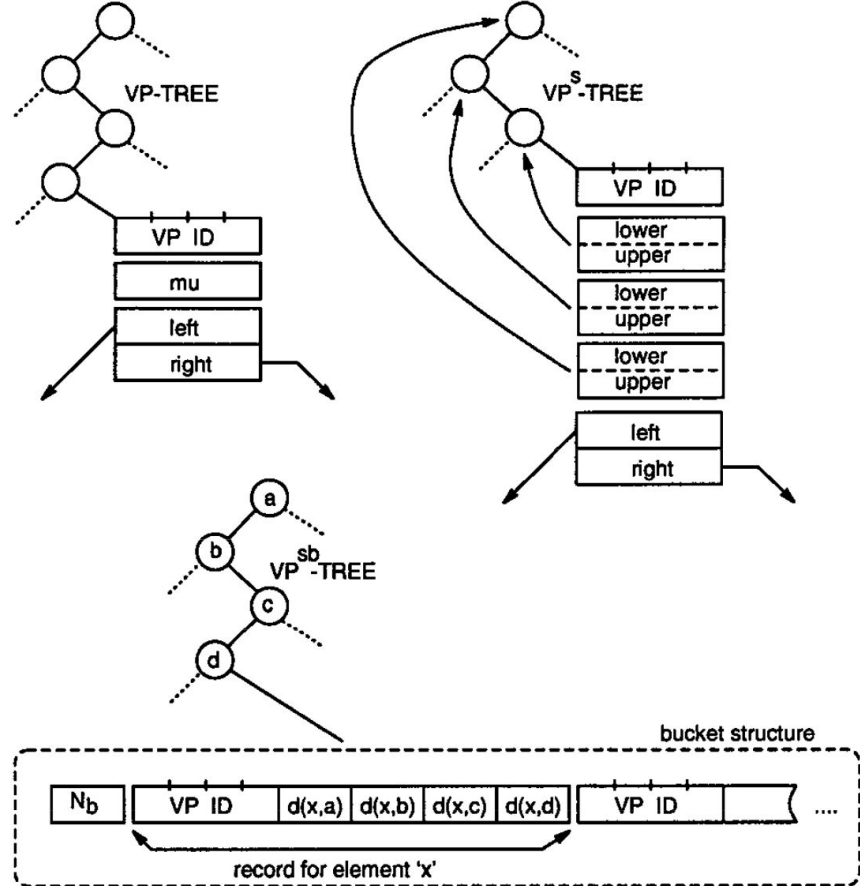


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# Performance comparison with classical metric trees

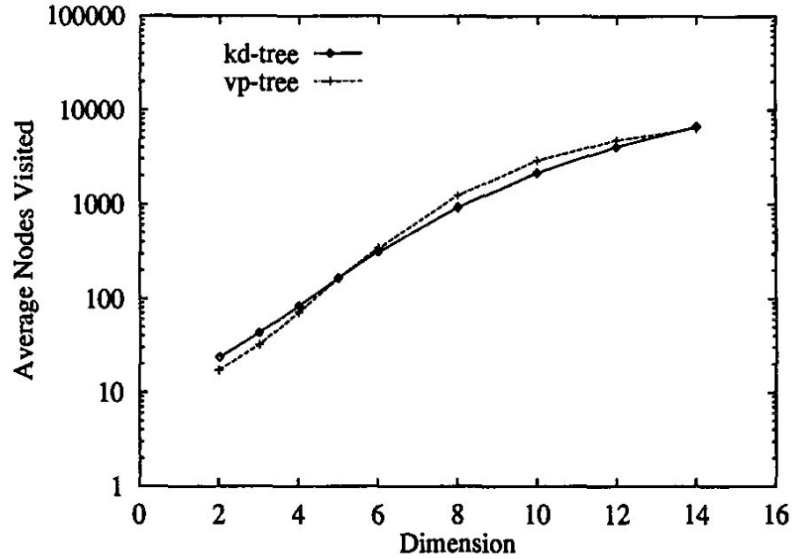


Figure 6: Search Cost vs. Dimension Comparison For  $L_2$  Metric and Based on Nodes Visited

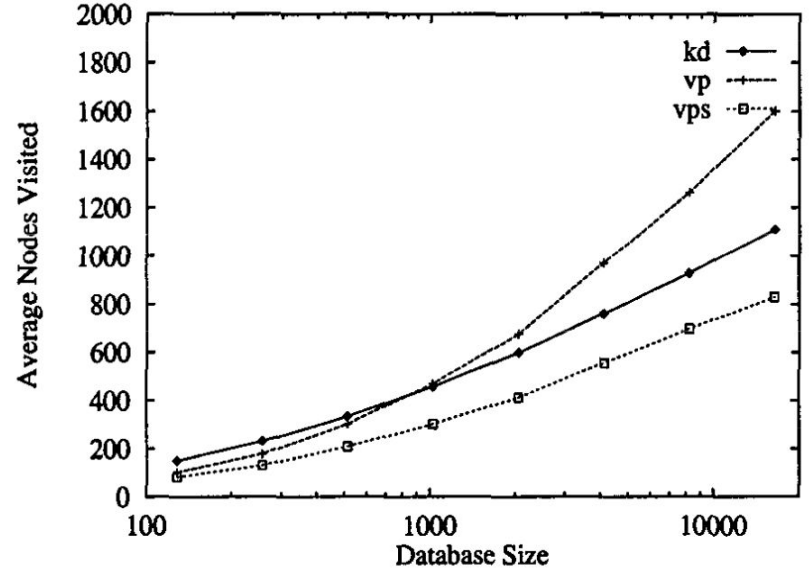


Figure 7: Search Cost vs. Database Size - Dimension 8

# Performance comparison with classical metric trees

- “What Is a Good Nearest Neighbors Algorithm for Finding Similar Patches in Images?” (Neeraj et al., 2008)

**Table 2.** Summary of results. The *vp*-tree performs well in all respects.

Method	Construction Performance	$\epsilon$ -NN Search Performance	$k$ -NN Search Performance
<i>k</i> d-Tree	Excellent	Poor	Poor
PCA Tree	Poor	Fair	Fair
Ball Tree	Fair	Excellent	Excellent
<i>k</i> -Means	Poor	Good	Good
<i>vp</i> -Tree	Excellent	Excellent	Excellent

**Merci pour votre attention !**